C-BGP

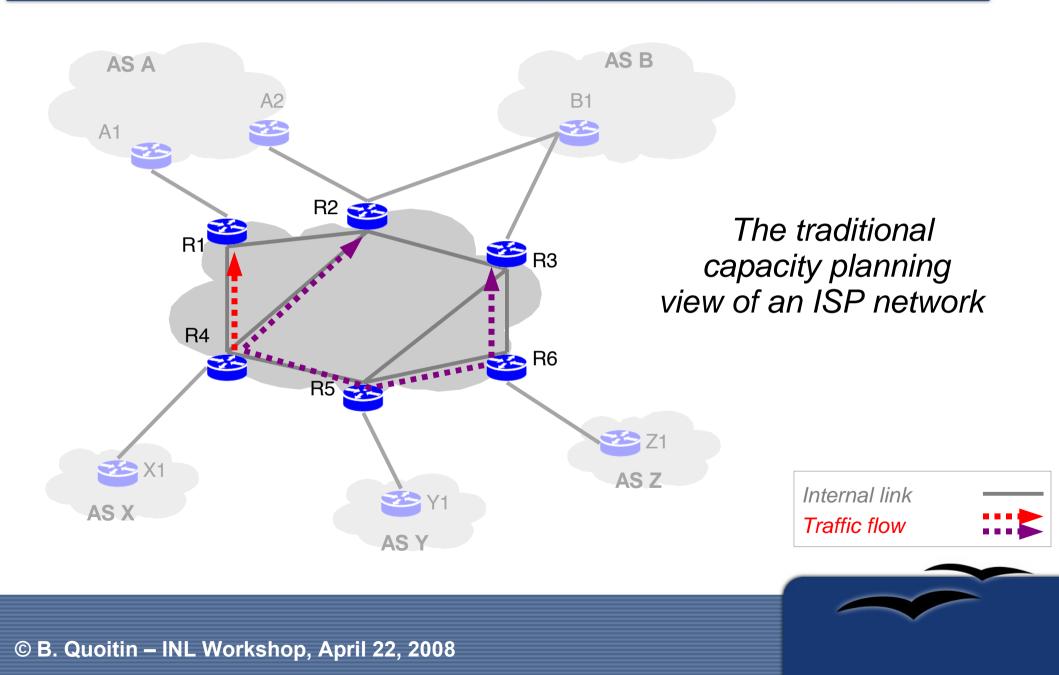
Modeling the Routing of an ISP

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ISP Model



Can't answer questions such as...

What would happen to my interdomain traffic if...

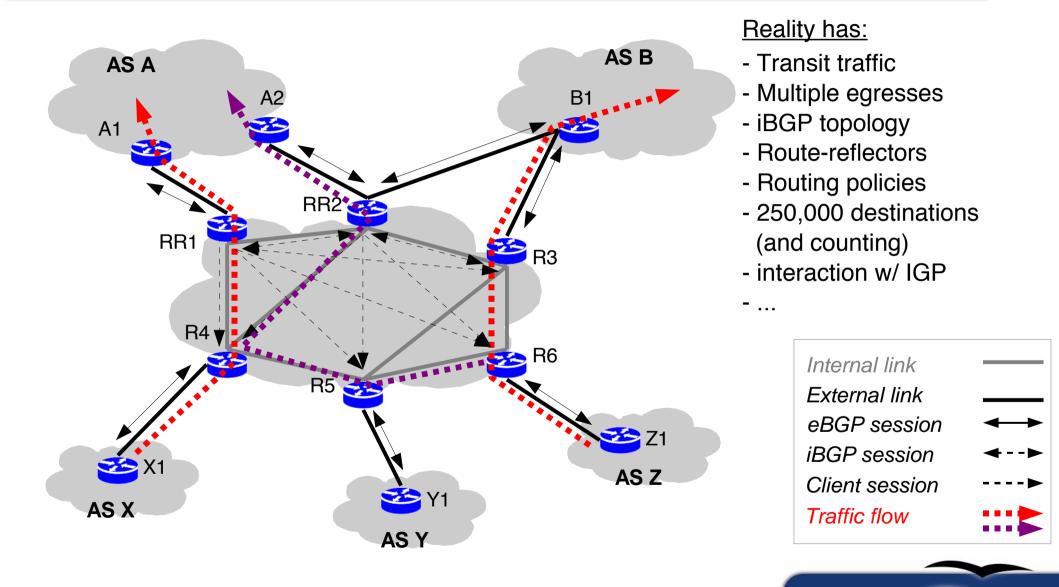
- a link is failing ?
- a router is under maintenance ?
- a BGP peering is being shutdown ?
- a new route filtering policy is planned ?
- a new peering is established at an IXP ?

How would I optimize my interdomain routing for...

- performance ?
- cost ?
- reliability ?



ISP Model



Agenda

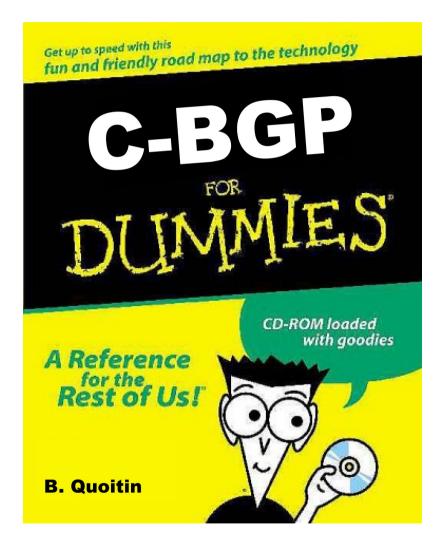
- C-BGP:
 - a network topology / config / routing DB
 - a BGP routing solver

- Case studies
 - Scenario 1: peering placement
 - Scenario 2: all single-link failures

I. C-BGP



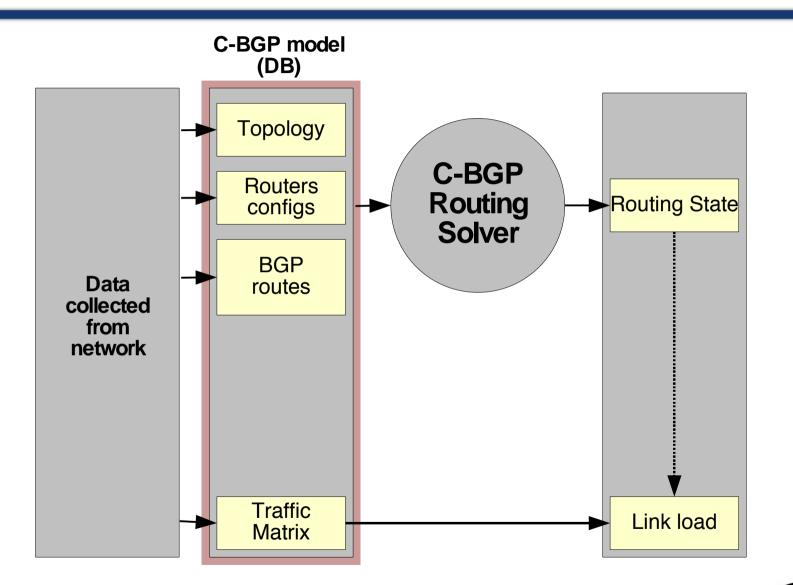
C-BGP



- Network topology / configuration DB
- Scriptable BGP Routing solver for large-scale networks
- developped by INL@UCLouvain
- supported by:



C-BGP Database



C-BGP Database

- Network topology database (DB)
 - Layer-3 model:
 - nodes: routers and LANs (pseudo-nodes)
 - links: IP links (need to be numbered)
 - link attributes: latency, bandwidth, load
 - Additional information:
 - static routes
 - IP tunnels
 - Large-scale topologies
 - tried with > 10000 nodes / 60000 links



DB: importing...

Build project from real router/network data



"show isis database extensive"

"show bgp neighbors"

JunOS IS-IS LSDB (TXT / XML)

- JunOS OSPF LSDB (TXT)
- IOS OSPF LSDB
- IOS IS-IS LSDB
- JunOS BGP neighbors (XML)
- IOS Running config
- JunOS Running config

_ (1)

🕞 Abilene

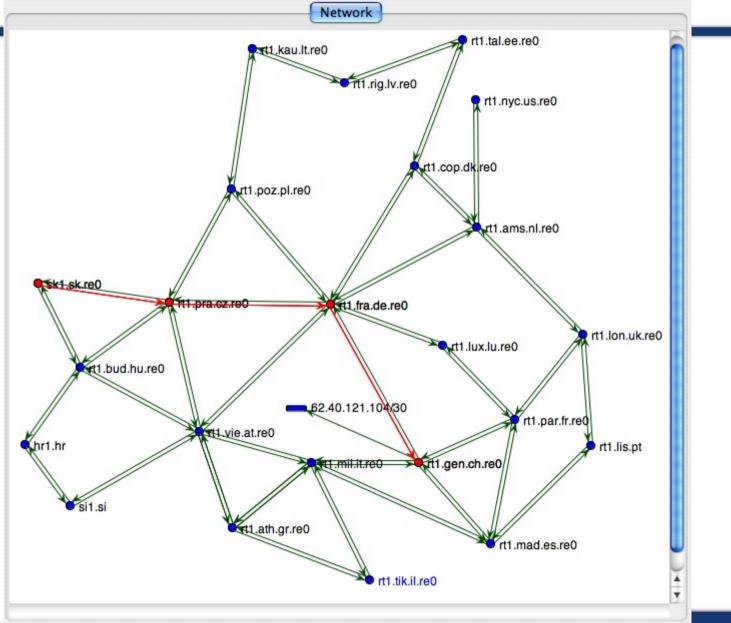
Project Manager

- Failers
 JunOS_ISIS (extensive,XML)
 JunOS BGP (XML)
 /Users/bqu/spinnet/workspace
 Data Groups
 - 🕨 📋 Group: default

⁽¹⁾ other vendor formats can be easily added



DB: network visualization

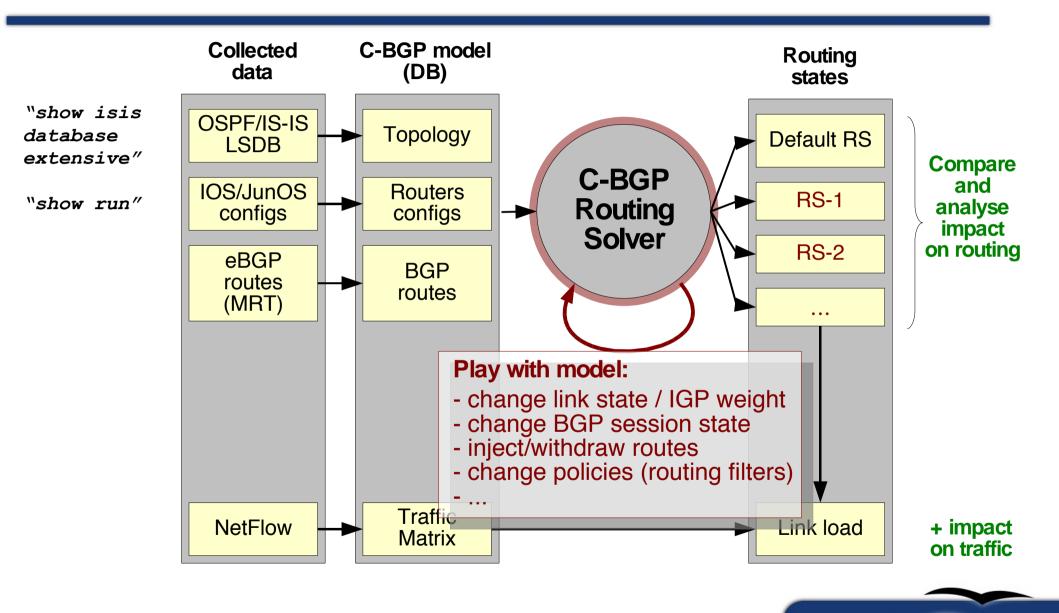


DB: network queries (CLI)

000		Console	2	
<pre>cbgp> show cbgp vers: libgds ver cbgp> net cbgp-node ptp ptp ptp ptp lo ptmp cbgp-node cbgp> net</pre>	<pre>ion: 1.5.0 [21 rsion: 1.4.5 node 62.40.11 > show ifaces 62.40.112.21/ 62.40.112.25/ 62.40.112.29/ 62.40.112.33/ 62.40.114.3/3 62.40.121.105 node 62.40.10 > traceroute 6 62.40.96.41</pre>	ib] [jni] [bgpdu 4.3 /30 /30 /30 /30 /30 /30 /2.37 (2.37 (2.40.114.3 (62.40.114.5) (62.40.114.7)	icmp error	(time-exceeded) (time-exceeded)
cbgp-node	:>			



C-BGP Routing Solver

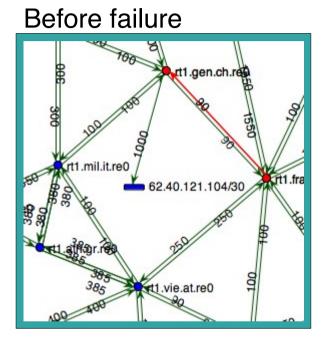


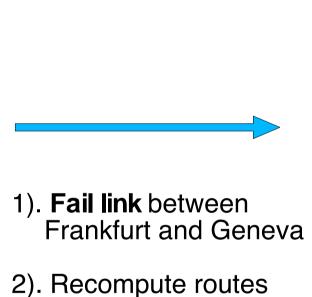
Routing Solver

- IGP model
 - static, centralized computation
 - support for ECMP, single area (currently)
- BGP model
 - compute steady-state outcome of BGP convergence
 - full decision process
 - versatile route filters
 - iBGP hierarchy (route-reflectors)
 - reads BGP table dumps and BGP message traces in MRT format (draft-ietf-grow-mrt-07.txt)



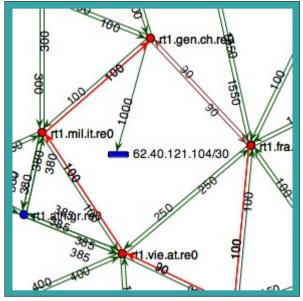
What-if scenarios





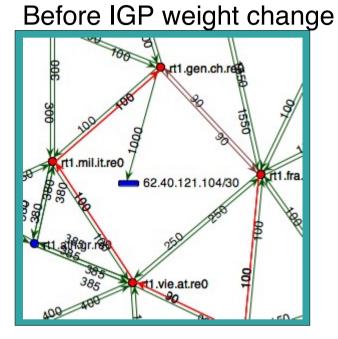
3). Trace route

After failure



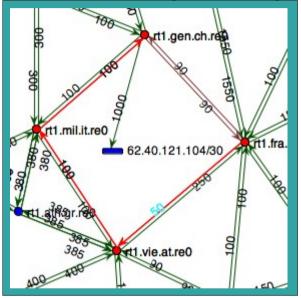


What-if scenarios



- 1). Change **IGP weight** of link between Frankfurt and Vienna
- 2). Recompute routes
- 3). Trace route

After IGP weight change



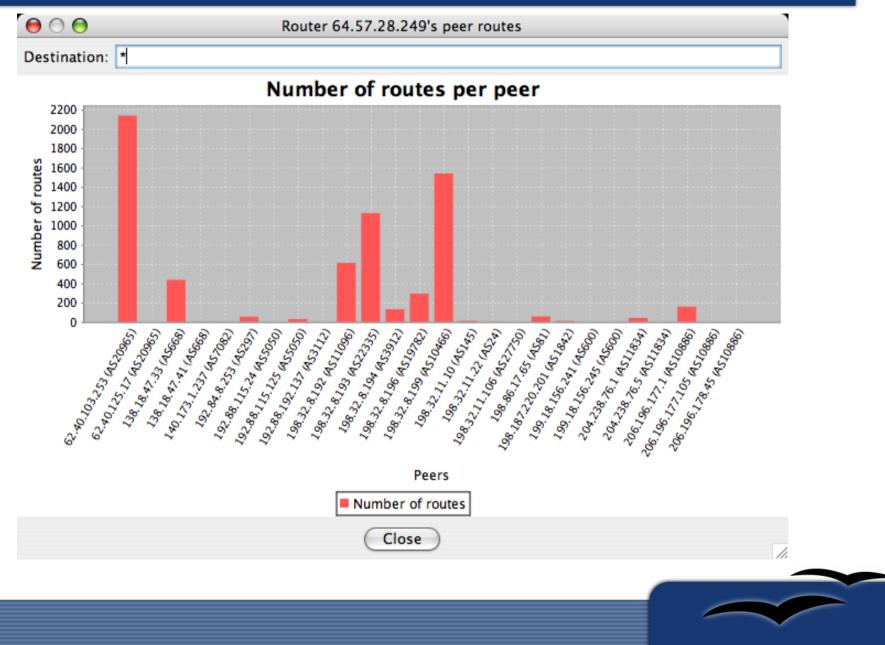


Routing State Queries

BGP domains	Inf	o Peers	Netw	orks R	IB	
AS 11537						
64.57.28.241 (CHIC-re0) Destination: *						
64.57.28.242 (NEWY-re0)	Selection size: 5145 rd	ute(s)				
64.57.28.243 (ATLA-re1)						
64.57.28.245 (KANS-re0)	E F Prefix	Next-Ho	Local-	MED Ori		
64.57.28.246 (SALT-re1)	8.6.244 1		200		11096 6	
64.57.28.248 (LOSA-re0)	 ✓ ✓ Ø 8.10.208 1 Ø 9.4.0.0/16 6 		260		10466 3	
	✓ ✓ 9.4.0.0/16 6 ✓ ✓ 12.0.48 1		100		20965 559	
64.57.28.249 (WASH-re0)			200		10578 1	
198.32.8.195 (HSTNng-re0)	 ✓ ✓ ✓ 12.6.208 1 ✓ ✓ 12.107.2 1 		200 200		10578 1 81 22753	
198.32.8.198 (LOSAng-re1)	▼ ▼ 12.107.2 1		260		10466 1	
198.32.8.200 (STTLng-re0)	▼ ▼ 12.151.0 1		260		10466 1	
198.32.8.201 (SNVAng-re1)	 ✓ ✓ ✓ 12.151.0 1 ✓ ✓ 12.151.1 1 		260		10466 1	
198.32.8.203 (ATLA-m5)	▼ ▼ □ 12.161.8 1		260		10466 88	
	▼ ▼ 12.161.8 1 ▼ ▼ 12.174.2 1		200		5661 21	
	V V 18.3.4.0/ 1		200		10578 3	1153
		00.00	222.14	4 50 0/2		
	Destination			4.59.0/2	9	
	Best Feasible		true true			
	Internal		false			
	Next-hop			2.8.199		
	AS-Path			6 13778		
	Local-Pref		260	0 15//0		
	MED		0			
	Communities		1153	7:2000 11	1537:950 1	1537:260



Routing State Queries



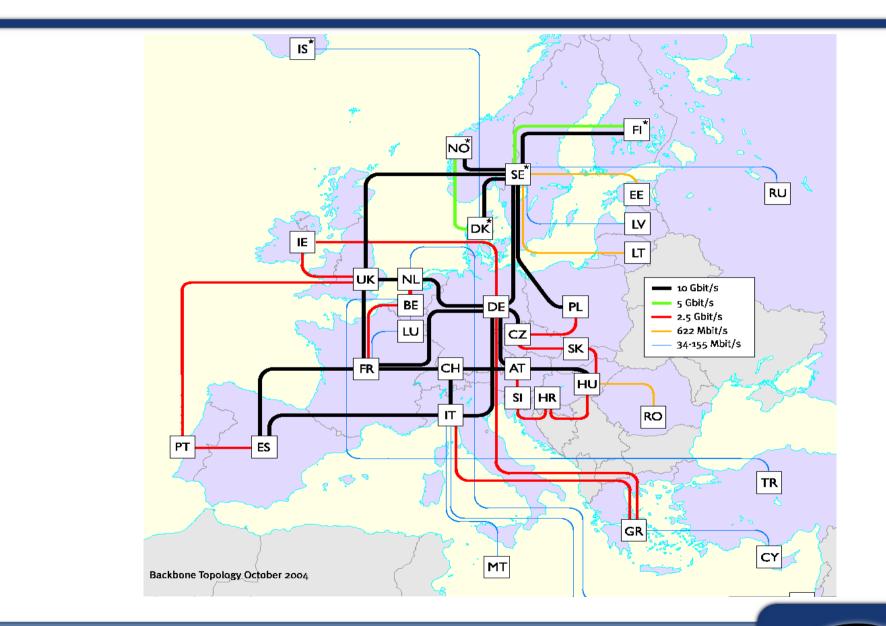
II. Case studies



Anonymous French Tier-1



Case study: GEANT (AS20965)



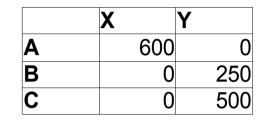
Case study: GEANT (AS20965)

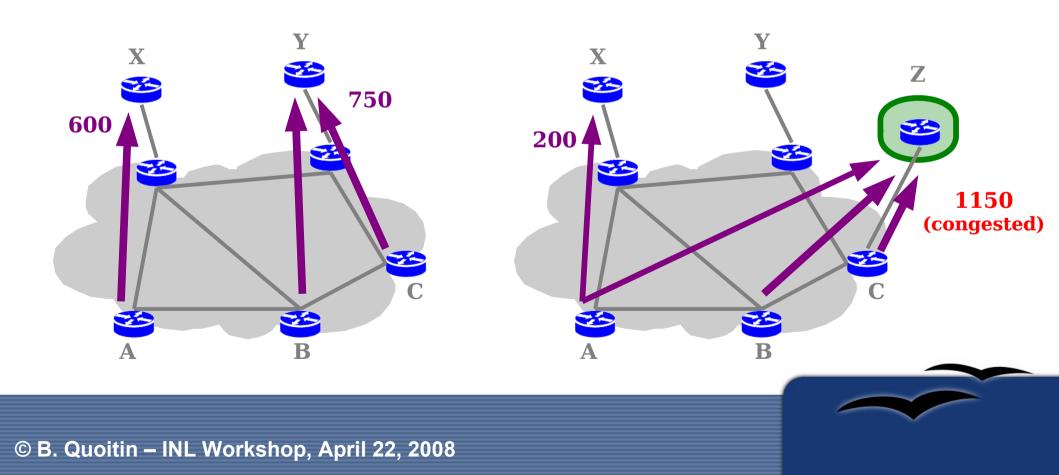
- Topology
 - Obtained from IS-IS trace, cross-checked with map
 - 23 nodes, 38 core links, 53 edge links (6 with upstreams)

Routing data

- Collected using Zebra in the iBGP (only best eBGP)
- 640,897 eBGP routes
 - 150,071 prefixes (clustered in 406 groups)
- Traffic data
 - NetFlow collected on all external interfaces
 - Sampling rate: 1/1000
 - About 150 GB per month
 - Src. / dst. aggregated in /24 prefixes

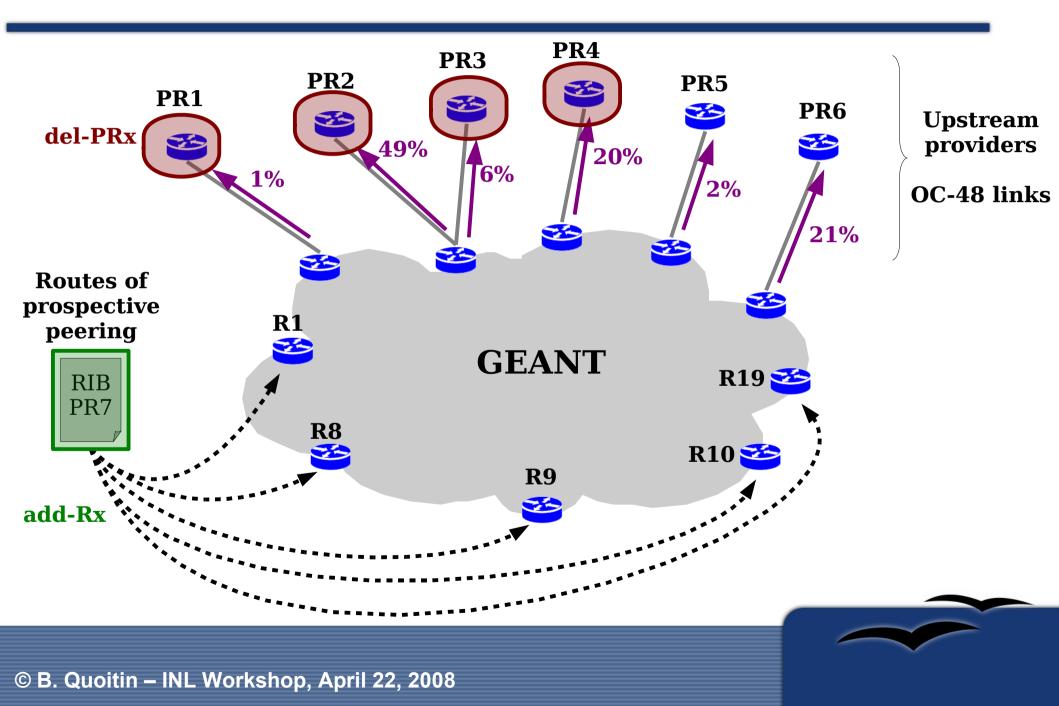
- Example
 - 2 upstream providers, 1Gbps links
 - Peer with new provider Z in C

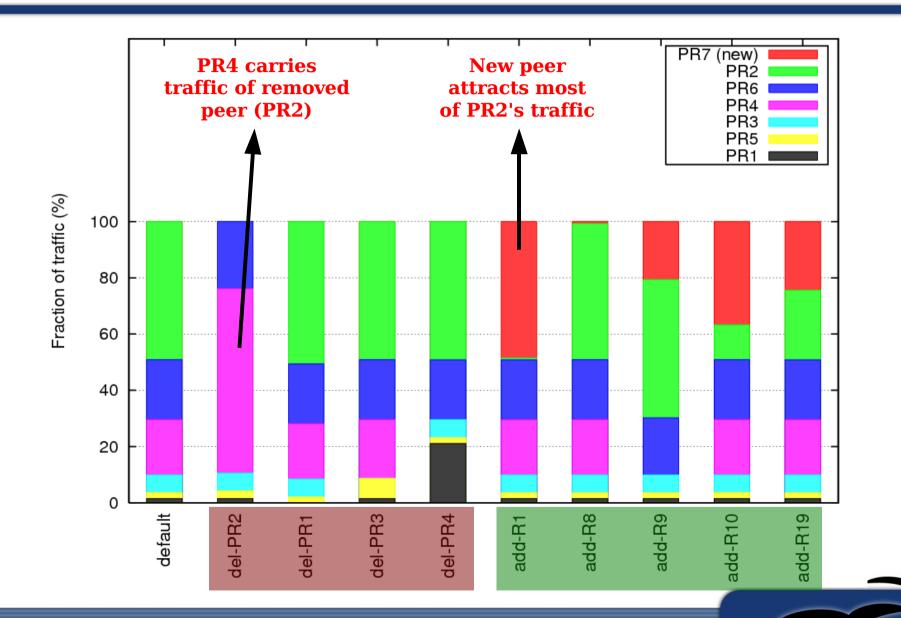




- Objective
 - Investigate addition/removal of peerings
 - Goal: better balance traffic load, reduce peering cost, ...
- Methodology
 - Scenario add-Rx
 - Consider a prospective peering **PR** (full RIB)
 - Inject routes of **PR** at router **Rx**
 - Scenario del-PRx
 - Remove the routes learned from an existing peer **PRx**
 - Metric
 - distribution of traffic among peering links (here: 6 most important links, OC-48 with upstream providers)

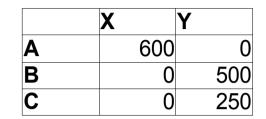


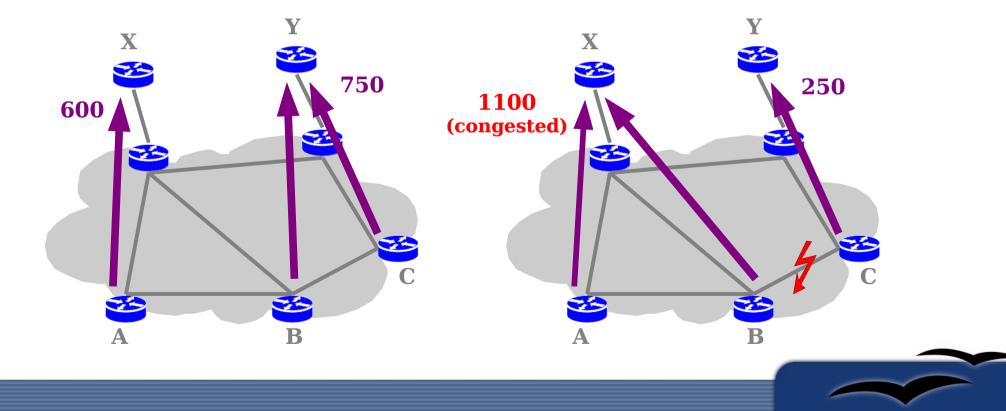




2nd Scenario: link failures

- Example
 - Traffic to upstream X and Y
 - 1 Gbps links
 - Internal link failure: $B \leftrightarrow C$

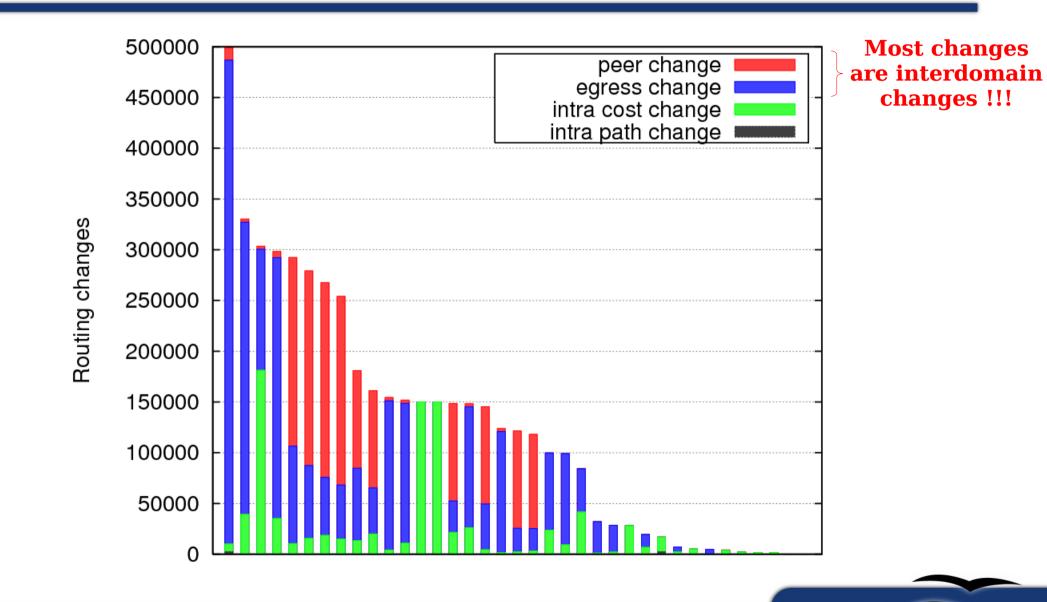




2nd Scenario: link failures

- Objectives
 - Study impact of single-link internal failures on routing
 - Consider all interdomain routes
- Methodology
 - Classification of routing changes
 - Prefix reachability
 - Peer change: neighbor AS has changed
 - Egress change: same AS, egress router changed
 - Intra cost change: same egress, IGP cost changed
 - Intra path change: same egress, same IGP cost, path changed (only when ECMP is allowed)

2nd Scenario: link failures



Conclusion

- Modeling the routing of an ISP is complex !
 - Many parameters and data sources are involved.
- Tools such as C-BGP
 - Helps understand & visualize routing protocol interaction (IGP / BGP) on large networks with many destinations
 - Useful to predict impact of events / choices on routing
 - Can be used as/with a capacity planning tool (if TM available) to predict the impact on link load



Thanks for your attention !

Visit our website: http://cbgp.info.ucl.ac.be

The C-BGP core is released under the LGPL license. The C-BGP core has been tested under the following platforms:







References

- Modeling the Routing of an ISP Network, B. Quoitin and S. Uhlig, IEEE Network, Vol 19(6), November 2005.
- Semi-automatic AS-wide converter for C-BGP, S. Tandel. Available from http://alumni.info.ucl.ac.be/standel/bgp-converter
- **Providing public intradomain traffic matrices to the research community**, S. Uhlig, B. Quoitin, S. Balon and J. Lepropre, ACM SIGCOMM Computer Communication Review, Vol 36(1), January 2006.
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